Specification of Source §2 Stepper—2022 edition

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Reduction

The *reducer* $\Rightarrow$ is a partial function from programs/statements/expressions to programs/statements/expressions (with slight abuse of notation via overloading) and $\Rightarrow^*$ is its reflexive transitive closure. A *reduction* is a sequence of programs $p_1 \Rightarrow \cdots \Rightarrow p_n$, where p_n is not reducible, i.e. there is no program q such that $p_n \Rightarrow q$. Here, the program p_n is called the *result of reducing* p_1.
If p_n, the result of reducing a program, is of the form v`;`, we call v the result of the program evaluation. If p_n is the empty sequence of statements, we declare the result of the program evaluation to be the value undefined. Reduction can get "stuck": the result can be a program that has neither of these two forms, which corresponds to a runtime error.
A *value* is a primitive number expression, primitive boolean expression, a primitive string expression, a function definition expression or a function declaration expression.
The *substitution* function $p[n := v]$ on programs/statements/expressions replaces every free occurrence of the name n in statement p by value v. Care must be taken to introduce and preserve co-references in this process; substitution can introduce cyclic references in the result of the substitution. For example, n may occur free in v, in which case every occurrence of n in p will be replaced by v such that n in v refers cyclically to the node at which the replacement happens.

Programs

Program-intro: In a sequence of statements, we can always reduce the first one that isn't a value statement.

$$\frac{\textit{statement} \Rightarrow \textit{statement}'}{[\textit{value}\texttt{;}]\ \ \textit{statement} \ldots \Rightarrow [\textit{value}\texttt{;}]\ \ \textit{statement}' \ldots}$$

Program-reduce: When the first two statements in a program are value statements, the first one can be removed.

$$\frac{\textit{v1}, \textit{v2} \text{ are values}}{\textit{v1}\texttt{;}\ \ \textit{v2}\texttt{;}\ \ \textit{statement} \ldots \Rightarrow \textit{v2}\texttt{;}\ \ \textit{statement} \ldots}$$

Eliminate-function-declaration: Function declarations as the first statement optionally after one value statement are substituted in the remaining statements.

$$\frac{f = \texttt{function}\ \textit{name}\ \texttt{(}\ \textit{parameters}\ \texttt{)}\ \textit{block}}{[\textit{value}\texttt{;}]\ \ f\ \textit{statement} \ldots\ \Rightarrow\ [\textit{value}\texttt{;}]\ \ \textit{statement} \ldots [\textit{name} := f]}$$

Eliminate-constant-declaration: Constant declarations as the first statement optionally after one value statement are substituted in the remaining statements.

$$\frac{c = \texttt{const}\ \textit{name}\ \texttt{=}\ v}{[\textit{value}\texttt{;}]\ \ c\ \textit{statement} \ldots\ \Rightarrow\ [\textit{value}\texttt{;}]\ \ \textit{statement} \ldots [\textit{name} := v]}$$

Statements: Expression statements

Expression-statement-reduce: An expression statement is reducible if its expression is reducible.

$$\frac{e \;\Rightarrow\; e'}{e\texttt{;} \;\Rightarrow\; e'\texttt{;}}$$

Statements: Constant declarations

Evaluate-constant-declaration: The right-hand expressions in constant declarations are evaluated.

$$\frac{\textit{expression} \;\Rightarrow\; \textit{expression}'}{\texttt{const}\ \textit{name}\ \texttt{=}\ \textit{expression} \;\Rightarrow\; \texttt{const}\ \textit{name}\ \texttt{=}\ \textit{expression}'}$$

Statements: Conditionals

Conditional-statement-predicate: A conditional statement is reducible if its predicate is reducible.

$$\frac{e \;\Rightarrow\; e'}{\texttt{if (}\ e\ \texttt{) \{}\ \cdots \texttt{\} else \{}\ \cdots \texttt{\}} \;\Rightarrow\; \texttt{if (}\ e'\ \texttt{) \{}\ \cdots \texttt{\} else \{}\ \cdots \texttt{\}}}$$

SICP-JS Exercise 4.8 specifies that a conditional statement where the taken branch is non-value-producing is value-producing, with its value being `undefined`.

Conditional-statement-consequent: Immediately within a program or block statement, a conditional statement whose predicate is true reduces to the consequent block.

$$\frac{}{\texttt{if (true) \{}\ \textit{program}_1 \texttt{\} else \{}\ \textit{program}_2 \texttt{\}} \;\Rightarrow\; \texttt{\{ undefined;}\ \ \textit{program}_1 \texttt{\}}}$$

Conditional-statement-alternative: Immediately within a program or block statement, a conditional statement whose predicate is false reduces to the alternative block.

$$\frac{}{\texttt{if (false) \{}\ \textit{program}_1 \texttt{\} else \{}\ \textit{program}_2 \texttt{\}} \;\Rightarrow\; \texttt{\{ undefined;}\ \ \textit{program}_2 \texttt{\}}}$$

We can remove the `undefined;` if the conditional statement is in a block expression without affecting the return value of the block expression. This is performed to reduce the surprise factor.

Conditional-statement-blockexpr-consequent: Immediately within a block expression, a conditional statement whose predicate is true reduces to the consequent block.

$$\frac{}{\texttt{\{}\ [\textit{value}\texttt{;}]\ \texttt{if (true) \{}\ \textit{program}_1 \texttt{\} else \{}\ \textit{program}_2 \texttt{\} \}} \;\Rightarrow\; \texttt{\{}[\textit{value}\texttt{;}]\ \texttt{\{}\ \textit{program}_1 \texttt{\}\}}}$$

Conditional-statement-blockexpr-alternative: Immediately within a block expression, a conditional statement whose predicate is false reduces to the alternative block.

$$\frac{}{\texttt{\{}\ [\textit{value}\texttt{;}]\ \texttt{if (false) \{}\ \textit{program}_1 \texttt{\} else \{}\ \textit{program}_2 \texttt{\} \}} \;\Rightarrow\; \texttt{\{}[\textit{value}\texttt{;}]\ \texttt{\{}\ \textit{program}_2 \texttt{\}\}}}$$

Statements: Blocks

Block-statement-intro: A block statement is reducible if its body is reducible.

$$\frac{\textit{program} \;\Rightarrow\; \textit{program}'}{\texttt{\{}\ \textit{program}\ \texttt{\}} \;\Rightarrow\; \texttt{\{}\ \textit{program}'\ \texttt{\}}}$$

Block-statement-single-reduce: A block statement consisting of a single value statement reduces to that value statement.

$$\frac{}{\texttt{\{}\textit{value}\texttt{;}\,\texttt{\}} \Rightarrow \textit{value}\texttt{;}}$$

Block-statement-empty-reduce: An empty block statement (or a non-value- producing one, which will reduce to the empty block statement) is non-value-producing, i.e. reduces to ε.

$$\frac{}{\texttt{\{\}} \Rightarrow \varepsilon}$$

Expresssions: Blocks

Block-expression-intro: A block expression is reducible if its body is reducible.

$$\frac{\textit{program} \;\Rightarrow\; \textit{program}'}{\texttt{\{}\;\textit{program}\;\texttt{\}} \;\Rightarrow\; \texttt{\{}\;\textit{program}'\;\texttt{\}}}$$

Block-expression-single-reduce: A block expression consisting of a single value statement reduces to **`undefined`**.

$$\frac{}{\texttt{\{}\textit{value}\texttt{;}\,\texttt{\}} \Rightarrow \textbf{\texttt{undefined}}}$$

Block-expression-empty-reduce: An empty block expression reduces to **`undefined`**.

$$\frac{}{\texttt{\{\}} \Rightarrow \textbf{\texttt{undefined}}}$$

Block-expression-return-reduce-1: In a block expression starting with a value statement and then a return statement, the value statement can be removed.

$$\frac{}{\texttt{\{}\textit{value}\texttt{;}\;\; \textbf{\texttt{return}}\;\; \textit{return-expression}\texttt{;}\;\; \textit{statement}\ldots\texttt{\}} \;\Rightarrow\; \texttt{\{}\textbf{\texttt{return}}\;\; \textit{return-expression}\texttt{;}\;\; \textit{statement}\ldots\texttt{\}}}$$

Block-expression-return-reduce-2: A block expression starting with a return statement reduces to the expression of the return statement.

$$\frac{}{\texttt{\{}\textbf{\texttt{return}}\;\; \textit{return-expression}\texttt{;}\;\; \textit{statement}\ldots\texttt{\}} \Rightarrow \textit{return-expression}}$$

Block expressions are currently used only as expanded forms of functions.

Expressions: Binary operators

Left-binary-reduce: An expression with binary operator can be reduced if its left sub-expression can be reduced.

$$\frac{e_1 \;\Rightarrow\; e_1'}{e_1\; \textit{binary-operator}\; e_2 \;\Rightarrow\; e_1'\; \textit{binary-operator}\; e_2}$$

And-shortcut-false: An expression with binary operator **`&&`** whose left sub-expression is **`false`** can be reduced to **`false`**.

$$\frac{}{\textbf{\texttt{false}}\;\textbf{\texttt{\&\&}}\; e \;\Rightarrow\; \textbf{\texttt{false}}}$$

And-shortcut-true: An expression with binary operator **`&&`** whose left sub-expression is **`true`** can be reduced to the right sub-expression.

$$\frac{}{\textbf{\texttt{true}}\;\textbf{\texttt{\&\&}}\; e \;\Rightarrow\; e}$$

Or-shortcut-true: An expression with binary operator `||` whose left sub-expression is `true` can be reduced to `true`.

$$\frac{}{\texttt{true || } e \;\Rightarrow\; \texttt{true}}$$

Or-shortcut-false: An expression with binary operator `||` whose left sub-expression is `false` can be reduced to the right sub-expression.

$$\frac{}{\texttt{false || } e \;\Rightarrow\; e}$$

Right-binary-reduce: An expression with binary operator can be reduced if its left sub-expression is a value and its right sub-expression can be reduced.

$$\frac{e_2 \;\Rightarrow\; e_2', \text{and } \textit{binary-operator} \text{is not } \texttt{\&\&} \text{ or } \texttt{||}}{v \textit{ binary-operator } e_2 \;\Rightarrow\; v \textit{ binary-operator } e_2'}$$

Prim-binary-reduce: An expression with binary operator can be reduced if its left and right sub-expressions are values and the corresponding function is defined for those values.

$$\frac{v \text{ is result of } v_1 \textit{ binary-operator } v_2}{v_1 \textit{ binary-operator } v_2 \;\Rightarrow\; v}$$

Expressions: Unary operators

Unary-reduce: An expression with unary operator can be reduced if its sub-expression can be reduced.

$$\frac{e \;\Rightarrow\; e'}{\textit{unary-operator } e \;\Rightarrow\; \textit{unary-operator } e'}$$

Prim-unary-reduce: An expression with unary operator can be reduced if its sub-expression is a value and the corresponding function is defined for that value.

$$\frac{v' \text{ is result of } \textit{unary-operator } v}{\textit{unary-operator } v \;\Rightarrow\; v'}$$

Expressions: conditionals

Conditional-predicate-reduce: A conditional expression can be reduced, if its predicate can be reduced.

$$\frac{e_1 \;\Rightarrow\; e_1'}{e_1 \texttt{ ? } e_2 \texttt{ : } e_3 \;\Rightarrow\; e_1' \texttt{ ? } e_2 \texttt{ : } e_3}$$

Conditional-true-reduce: A conditional expression whose predicate is the value `true` can be reduced to its consequent expression.

$$\frac{}{\texttt{true ? } e_1 \texttt{ : } e_2 \;\Rightarrow\; e_1}$$

Conditional-false-reduce: A conditional expression whose predicate is the value `false` can be reduced to its alternative expression.

$$\frac{}{\texttt{false ? } e_1 \texttt{ : } e_2 \;\Rightarrow\; e_2}$$

Expressions: function application

Application-functor-reduce: A function application can be reduced if its functor expression can be reduced.

$$\frac{e \ \Rightarrow \ e'}{e \ \texttt{(} \ \textit{expressions} \ \texttt{)} \ \Rightarrow \ e' \ \texttt{(} \ \textit{expressions} \ \texttt{)}}$$

Application-argument-reduce: A function application can be reduced if one of its argument expressions can be reduced and all preceding arguments are values.

$$\frac{e \ \Rightarrow \ e'}{v \ \texttt{(} \ v_1 \ldots v_i \ e \ldots \ \texttt{)} \ \Rightarrow \ v \ \texttt{(} \ v_1 \ldots v_i \ e' \ldots \ \texttt{)}}$$

Function-declaration-application-reduce: The application of a function declaration can be reduced, if all arguments are values.

$$\frac{f = \texttt{function} \ n \ \texttt{(} \ x_1 \ldots x_n \ \texttt{)} \ \textit{block}}{f \ \texttt{(} \ v_1 \ldots v_n \ \texttt{)} \ \Rightarrow \ \textit{block}[x_1 := v_1] \ldots [x_n := v_n][n := f]}$$

Function-definition-application-reduce: The application of a function definition can be reduced, if all arguments are values.

$$\frac{f = \texttt{(} \ x_1 \ldots x_n \ \texttt{)} \ \texttt{=>} \ b, \text{ where } b \text{ is an expression or block}}{f \ \texttt{(} \ v_1 \ldots v_n \ \texttt{)} \ \Rightarrow \ b[x_1 := v_1] \ldots [x_n := v_n]}$$

Substitution

Identifier: An identifier with the same name as x is substituted with e_x.

$$\frac{}{x[x := e_x] \ = \ e_x}$$

$$\frac{\textit{name} \ \neq \ x}{\textit{name}[x := e_x] \ = \ \textit{name}}$$

Expression statement: All occurrences of x in e are substituted with e_x.

$$\frac{}{e\texttt{;}\,[x := e_x] \ = \ e\,[x := e_x]\texttt{;}}$$

Binary expression: All occurrences of x in the operands are substituted with e_x.

$$\frac{}{(e_1 \ \textit{binary-operator} \ e_2)[x := e_x] \ = \ e_1[x := e_x] \ \textit{binary-operator} \ e_2[x := e_x]}$$

Unary expression: All occurrences of x in the operand are substituted with e_x.

$$\frac{}{(\textit{unary-operator} \ e)[x := e_x] \ = \ \textit{unary-operator} \ e[x := e_x]}$$

Conditional expression: All occurrences of x in the operands are substituted with e_x.

$$\frac{}{(e_1 \ \texttt{?} \ e_2 \ \texttt{:} \ e_3)[x := e_x] \ = \ e_1[x := e_x] \ \texttt{?} \ e_2[x := e_x] \ \texttt{:} \ e_3[x := e_x]}$$

Logical expression: All occurrences of x in the operands are substituted with e_x.

$$\frac{}{(e_1 \ \texttt{||} \ e_2)[x := e_x] \ = \ e_1[x := e_x] \ \texttt{||} \ e_2[x := e_x]}$$

$$\frac{}{(e_1 \ \texttt{\&\&} \ e_2)[x := e_x] \ = \ e_1[x := e_x] \ \texttt{\&\&} \ e_2[x := e_x]}$$

Call expression: All occurrences of x in the arguments and the function expression of the application e are substituted with e_x.

$$\frac{}{(e \ \texttt{(} \ x_1 \ldots x_n \ \texttt{)})[x := e_x] \ = \ e[x := e_x] \ \texttt{(} \ x_1[x := e_x] \ldots x_n[x := e_x] \ \texttt{)}}$$

Function declaration: All occurrences of x in the body of a function are substituted with e_x under given circumstances.

1. Function declaration where x has the same name as a parameter.

$$\frac{\exists\, i \in \{1, \cdots, n\} \text{ s.t. } x \ = \ x_i}{(\texttt{function} \ \textit{name} \ \texttt{(} \ x_1 \ldots x_n \ \texttt{)} \ \textit{block})[x := e_x] \ = \ \texttt{function} \ \textit{name} \ \texttt{(} \ x_1 \ldots x_n \ \texttt{)} \ \textit{block}}$$

2. Function declaration where x does not have the same name as a parameter.

(a) No parameter of the function occurs free in e_x.

$$\frac{\forall\, i \in \{1, \cdots, n\} \text{ s.t. } x \;\neq\; x_i,\;\; \forall\, j \in \{1, \cdots, n\} \text{ s.t. } x_j \text{ does not occur free in } e_x}{\begin{array}{c}(\texttt{function}\ \textit{name}\ \texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \textit{block})[x := e_x] \\ = \\ \texttt{function}\ \textit{name}\ \texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \textit{block}[x := e_x]\end{array}}$$

(b) A parameter of the function occurs free in e_x.

$$\frac{\forall\, i \in \{1, \cdots, n\} \text{ s.t. } x \;\neq\; x_i,\;\; \exists\, j \in \{1, \cdots, n\} \text{ s.t. } x_j \text{ occurs free in } e_x}{\begin{array}{c}(\texttt{function}\ \textit{name}\ \texttt{(}\ x_1 \ldots x_j \ldots x_n\ \texttt{)}\ \textit{block})[x := e_x] \\ = \\ (\texttt{function}\ \textit{name}\ \texttt{(}\ x_1 \ldots y \ldots x_n\ \texttt{)}\ \textit{block}[x_j := y])[x := e_x]\end{array}}$$

Substitution is applied to the whole expression again as to recursively detect and rename all parameters of the function declaration that clash with variables that occur free in e_x, at which point (i) takes place. Note that the name y is not declared in, nor occurs in *block* and e_x.

Lambda expression: All occurrences of x in the body of a lambda expression are substituted with e_x under given circumstances.

1. Lambda expression where x has the same name as a parameter.

$$\frac{\exists\, i \in \{1, \cdots, n\} \text{ s.t. } x \;=\; x_i}{(\texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block})[x := e_x] \;=\; \texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block}}$$

2. Lambda expression where x does not have the same name as a parameter.

(a) No parameter of the lambda expression occurs free in e_x.

$$\frac{\forall\, i \in \{1, \cdots, n\} \text{ s.t. } x \;\neq\; x_i,\;\; \forall\, j \in \{1, \cdots, n\} \text{ s.t. } x_j \text{ does not occur free in } e_x}{(\texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block})[x := e_x] \;=\; \texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block}[x := e_x]}$$

(b) A parameter of the lambda expression occurs free in e_x.

$$\frac{\forall\, i \in \{1, \cdots, n\} \text{ s.t. } x \;\neq\; x_i,\;\; \exists\, j \in \{1, \cdots, n\} \text{ s.t. } x_j \text{ occurs free in } e_x}{(\texttt{(}\ x_1 \ldots x_j \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block})[x := e_x] \;=\; (\texttt{(}\ x_1 \ldots y \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block}[x_j := y])[x := e_x]}$$

Substitution is applied to the whole expression again as to recursively detect and rename all parameters of the lambda expression that clash with variables that occur free in e_x, at which point (i) takes place. Note that the name y is not declared in, nor occurs in *block* and e_x.

Block expression: All occurrences of x in the statements of a block expression are substituted with e_x under given circumstances.

1. Block expression in which x is declared.

$$\frac{x \text{ is declared in } \textit{block}}{\textit{block}[x := e_x] \;=\; \textit{block}}$$

2. Block expression in which x is not declared.

   (a) No names declared in the block occurs free in e_x.

$$\frac{x \text{ is not declared in } \textit{block},\;\; \textit{name} \text{ declared in } \textit{block} \text{ does not occur free in } e_x}{\textit{block}[x := e_x] \;=\; [\textit{block}[0][x := e_x],\; \ldots,\; \textit{block}[n][x := e_x]]}$$

   (b) A name declared in the block occurs free in e_x.

$$\frac{x \text{ is not declared in } \textit{block},\;\; \textit{name} \text{ declared in } \textit{block} \text{ occurs free in } e_x}{\textit{block}[x := e_x] \;=\; [\textit{block}[0][\textit{name} := y],\; \ldots,\; \textit{block}[n][\textit{name} := y]][x := e_x]}$$

   Substitution is applied to the whole expression again as to recursively detect and rename all declared names of the block expression that clash with variables that occur free in e_x, at which point `(i)` takes place. Note that the name y is not declared in, nor occurs in *block* and e_x.

Variable declaration: All occurrences of x in the declarators of a variable declaration are substituted with e_x.

$$\frac{}{\textit{declarations}[x := e_x] \;=\; [\textit{declarations}[0][x := e_x] \ldots \textit{declarations}[n][x := e_x]]}$$

Return statement: All occurrences of x in the expression that is to be returned are substituted with e_x.

$$\frac{}{(\texttt{return}\; e\texttt{;})[x := e_x] \;=\; \texttt{return}\; e[x := e_x]\texttt{;}}$$

Conditional statement: All occurrences of x in the condition, consequent, and alternative expressions of a conditional statement are substituted with e_x.

$$\frac{}{(\texttt{if (}\; e \;\texttt{)}\; \textit{block}\; \texttt{else}\; \textit{block})[x := e_x] \;=\; \texttt{if (}\; e[x := e_x] \;\texttt{)}\; \textit{block}[x := e_x]\; \texttt{else}\; \textit{block}[x := e_x]}$$

Array expression: All occurrences of x in the elements of an array are substituted with e_x.

$$\frac{}{[x_1,\; \ldots,\; x_n][x := e_x] \;=\; [x_1[x := e_x],\; \ldots,\; x_n[x := e_x]]}$$

Free names

Let $\triangleright$ be the relation that defines the set of free names of a given Source expression; the symbols p_1 and p_2 shall henceforth refer to unary and binary operations, respectively. That is, p_1 ranges over $\{$`!`$\}$ and p_2 ranges over $\{$`||`, `&&`, `+`, `-`, `*`, `/`, `===`, `>`, `<`$\}$.

Identifier:

$$\frac{}{x \ \triangleright\ \{x\}}$$

$$\frac{}{\textit{name} \ \triangleright\ \varnothing}$$

Boolean:

$$\frac{}{\texttt{true} \ \triangleright\ \varnothing}$$

$$\frac{}{\texttt{false} \ \triangleright\ \varnothing}$$

Expression statement:

$$\frac{e \ \triangleright\ S}{e\texttt{;} \ \triangleright\ S}$$

Unary expression:

$$\frac{e \ \triangleright\ S}{p_1(e) \ \triangleright\ S}$$

Binary expression:

$$\frac{e_1 \ \triangleright\ S_1,\ \ e_2 \ \triangleright\ S_2}{p_1(e_1, e_2) \ \triangleright\ S_1 \cup S_2}$$

Conditional expression:

$$\frac{e_1 \ \triangleright\ S_1,\ \ e_2 \ \triangleright\ S_2,\ \ e_3 \ \triangleright\ S_3}{e_1 \ \texttt{?}\ e_2 \ \texttt{:}\ e_3 \ \triangleright\ S_1 \cup S_2 \cup S_3}$$

Call expression:

$$\frac{e \ \triangleright\ S,\ \ e_k \ \triangleright\ T_k}{e(e_1,\ \ldots,\ e_n) \ \triangleright\ S \cup T_1 \cup \ldots \cup T_n}$$

Function declaration:

$$\frac{\textit{block} \ \triangleright\ S}{\texttt{function}\ \textit{name}\ \texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \textit{block} \ \triangleright\ S - \{x_1,\ \ldots,\ x_n\}}$$

Lambda expression:

$$\frac{\textit{block} \ \triangleright\ S}{\texttt{(}\ x_1 \ldots x_n\ \texttt{)}\ \texttt{=>}\ \textit{block} \ \triangleright\ S - \{x_1,\ \ldots,\ x_n\}}$$

Block expression:

$$\frac{\textit{block}[k] \ \triangleright \ S_k, \ T \text{ contains all names declared in } \textit{block}}{\textit{block} \ \triangleright \ (S_1 \cup \ldots \cup S_n) - T}$$

Constant declaration:

$$\frac{e \ \triangleright \ S}{\texttt{const } \textit{name} \texttt{ = } e\texttt{;} \ \triangleright \ S}$$

Return statement:

$$\frac{e \ \triangleright \ S}{\texttt{return } e\texttt{;} \ \triangleright \ S}$$

Conditional statement:

$$\frac{e \ \triangleright \ S, \ \textit{block}_1 \ \triangleright \ T_1, \ \textit{block}_2 \ \triangleright \ T_2}{\texttt{if (} e \texttt{) } \textit{block}_1 \texttt{ else } \textit{block}_2 \ \triangleright \ S \cup T_1 \cup T_2}$$